

# INTRODUCTION

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## Course Theme:

### Abstraction Is Good But Don't Forget Reality

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Most CS and CE courses emphasize abstraction

- Abstract data types
- Asymptotic analysis

These abstractions have limits

- Especially in the presence of bugs
- Need to understand details of underlying implementations

Useful outcomes

- Become more effective programmers
  - Able to find and eliminate bugs efficiently
  - Able to understand and tune for program performance
- Prepare for later "systems" classes
  - Compilers, Operating Systems, Networks, Computer Architecture, Embedded Systems

# Great Reality #1:

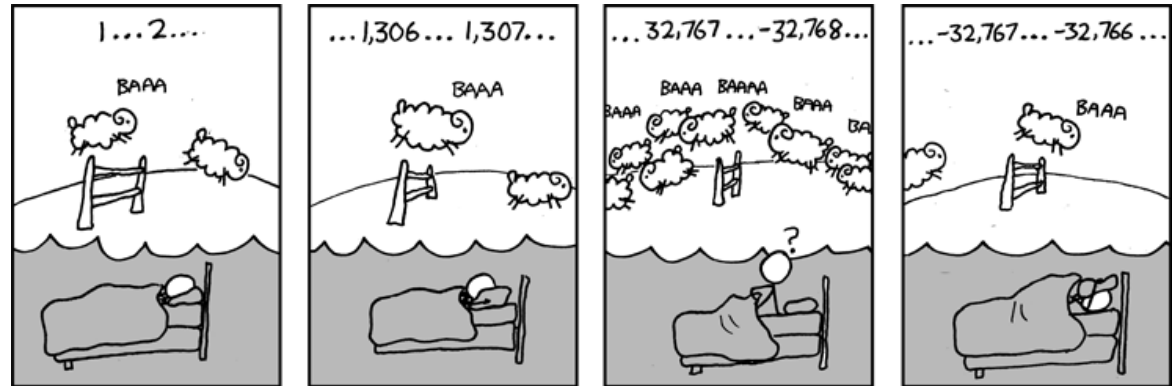
## Ints are not Integers, Floats are not Reals

Example 1: Is  $x^2 \geq 0$ ?

- Float's: Yes!

- Int's:

- $40000 * 40000 = 1600000000$
- $50000 * 50000 = -1794967296$  (overflow)



Example 2: Is  $(x + y) + z = x + (y + z)$  ?

- Unsigned & Signed Int's: Yes!
- Float's:
  - $(1e20 + -1e20) + 3.14 \rightarrow 3.14$
  - $1e20 + (-1e20 + 3.14) \rightarrow ??$

# Code Security Example

```
/* Kernel memory region holding user-accessible data */
#define KSIZE 1024
char kbuf[KSIZE];

/* Copy at most maxlen bytes from kernel region to user buffer */
int copy_from_kernel(void *user_dest, int maxlen) {
    /* Byte count len is minimum of buffer size and maxlen */
    int len = KSIZE < maxlen ? KSIZE : maxlen;
    memcpy(user_dest, kbuf, len);
    return len;
}
```

Similar to code found in FreeBSD's implementation of *getpeername()*

There are legions of smart people trying to find vulnerabilities in programs

# Typical Usage

```
/* Kernel memory region holding user-accessible data */
#define KSIZE 1024
char kbuf[KSIZE];

/* Copy at most maxlen bytes from kernel region to user buffer */
int copy_from_kernel(void *user_dest, int maxlen) {
    /* Byte count len is minimum of buffer size and maxlen */
    int len = KSIZE < maxlen ? KSIZE : maxlen;
    memcpy(user_dest, kbuf, len);
    return len;
}
```

```
#define MSIZE 528

void getstuff() {
    char mybuf[MSIZE];
    copy_from_kernel(mybuf, MSIZE);
    printf("%s\n", mybuf);
}
```

# Malicious Usage

```
/* Kernel memory region holding user-accessible data */
#define KSIZE 1024
char kbuf[KSIZE];

/* Copy at most maxlen bytes from kernel region to user buffer */
int copy_from_kernel(void *user_dest, int maxlen) {
    /* Byte count len is minimum of buffer size and maxlen */
    int len = KSIZE < maxlen ? KSIZE : maxlen;
    memcpy(user_dest, kbuf, len);
    return len;
}
```

```
#define MSIZE 528

void getstuff() {
    char mybuf[MSIZE];
    copy_from_kernel(mybuf, -MSIZE);
    . . .
}
```

# Computer Arithmetic

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Cannot assume all "usual" mathematical properties

- Due to finiteness of representations
- Integer operations satisfy "ring" properties
  - Commutativity, associativity, distributivity
- Floating point operations satisfy "ordering" properties
  - Monotonicity, values of signs

Observation

- Need to understand which abstractions apply in which contexts
- Important issues for compiler writers and serious application programmers

Chances are, you'll never write programs in assembly

- Compilers are much better & more patient than you are

But: Understanding assembly is key to machine-level execution model

- Behavior of programs in presence of bugs
  - High-level language models break down
- Tuning program performance
  - Understand optimizations done / not done by the compiler
  - Understanding sources of program inefficiency
- Implementing system software
  - Compiler has machine code as target
  - Operating systems must manage process state
- Creating / fighting malware



# Assembly Code Example

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## Time Stamp Counter

- Special 64-bit register in Intel-compatible machines
- Incremented every clock cycle
- Read with `rdtsc` instruction

## Application

- Measure time (in clock cycles) required by procedure

```
double t;  
start_counter();  
P();  
t = get_counter();  
printf("P required %f clock cycles\n", t);
```

# Code to Read Counter

Write small amount of assembly code using [GCC's asm facility](#)

Inserts assembly code into machine code generated by compiler

```
static unsigned cyc_hi = 0;
static unsigned cyc_lo = 0;

/* Set *hi and *lo to the high and low order bits
   of the cycle counter.
*/
void access_counter(unsigned *hi, unsigned *lo)
{
    asm("rdtsc; movl %%edx,%0; movl %%eax,%1"
        : "=r" (*hi), "=r" (*lo)
        :
        : "%edx", "%eax");
}
```

#### Memory referencing bugs especially pernicious

- Effects are distant in both time and space

#### Memory performance is not uniform

- `Cache` and `virtual memory` effects can greatly affect program performance
- Adapting program to characteristics of memory system can lead to major speed improvements

#### Memory is not unbounded

- It must be allocated and managed
- Many applications are memory dominated

# Memory Referencing Bug Example

```
double fun(int i)
{
    volatile double d[1] = {3.14};
    volatile long int a[2];
    a[i] = 1073741824; /* Possibly out of bounds */
    return d[0];
}
```

*volatile: Don't be optimized by compiler*

|        |                               |
|--------|-------------------------------|
| fun(0) | 3.14                          |
| fun(1) | 3.14                          |
| fun(2) | 3.1399998664856               |
| fun(3) | 2.00000061035156              |
| fun(4) | 3.14, then segmentation fault |

Result is architecture specific

# Memory Referencing Bug Example

```
double fun(int i)
{
    volatile double d[1] = {3.14};
    volatile long int a[2];
    a[i] = 1073741824; /* Possibly out of bounds */
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*volatile: Don't be optimized by compiler*

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| fun(0) | 3.14                          |
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Explanation:

|             |   |                                  |
|-------------|---|----------------------------------|
| Saved State | 4 | } Location accessed<br>by fun(i) |
| d7 ... d4   | 3 |                                  |
| d3 ... d0   | 2 |                                  |
| a[1]        | 1 |                                  |
| a[0]        | 0 |                                  |

# Memory Referencing Errors

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C and C++ do not provide any memory protection

- Out of bounds array references
- Invalid pointer values
- Abuses of malloc/free

Can lead to nasty bugs

- Whether or not bug has any effect depends on system and compiler
- Action at a distance
  - Corrupted object logically unrelated to one being accessed
  - Effect of bug may be first observed long after it is generated

How can I deal with this?

- Program in Java, Ruby or ML
- Understand what possible interactions may occur
- Use or develop tools to detect referencing errors (e.g. Valgrind)

# Memory System Performance Example

```
void copyij(int src[2048][2048],  
            int dst[2048][2048])  
{  
    int i,j;  
    for (i = 0; i < 2048; i++)  
        for (j = 0; j < 2048; j++)  
            dst[i][j] = src[i][j];  
}
```



```
void copyji(int src[2048][2048],  
            int dst[2048][2048])  
{  
    int i,j;  
    for (j = 0; j < 2048; j++)  
        for (i = 0; i < 2048; i++)  
            dst[i][j] = src[i][j];  
}
```

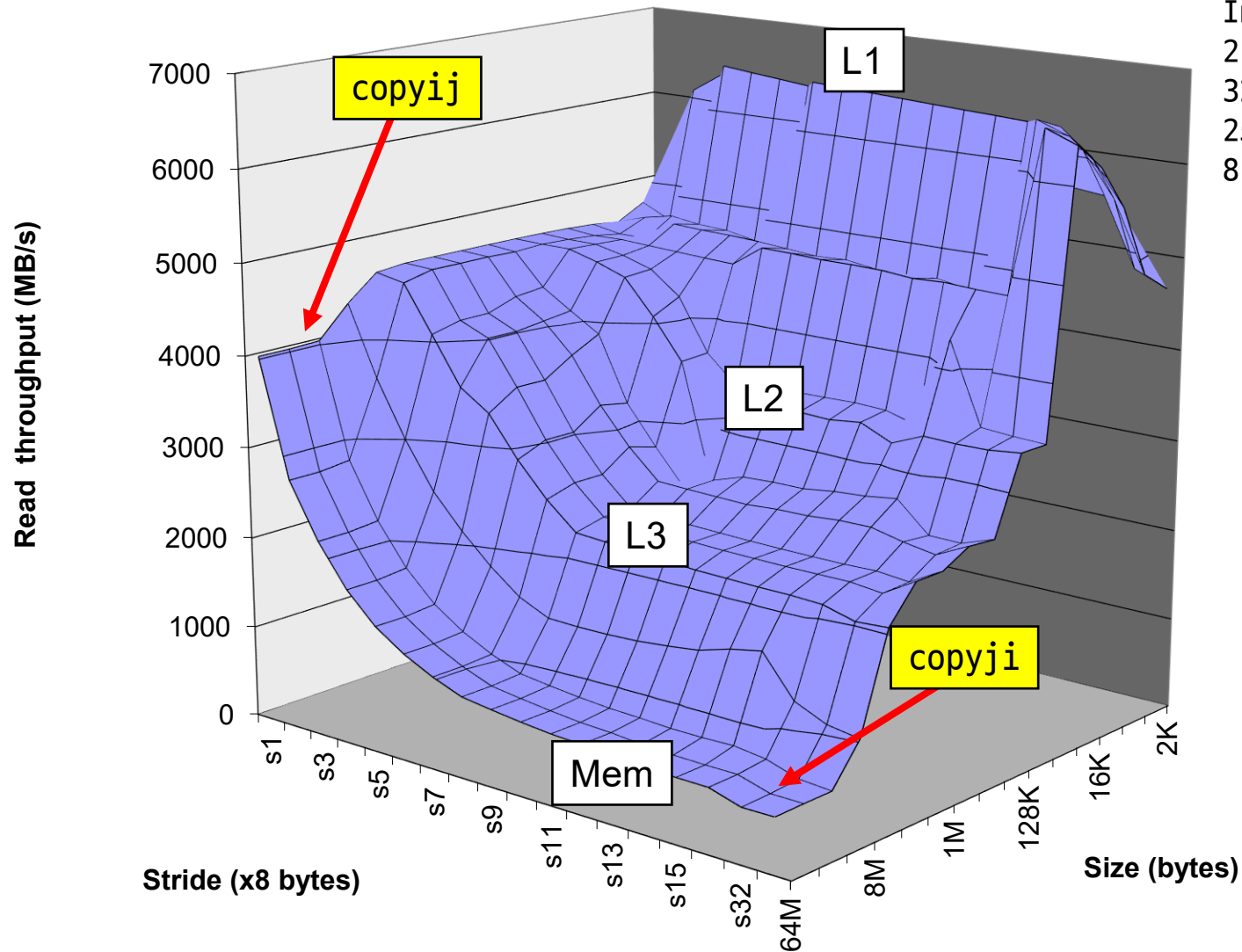
Need to understand hierarchical memory organization

Performance depends on access patterns

- Including how step through multi-dimensional array

# The Memory Mountain

Intel Core i7  
2.67 GHz  
32 KB L1 d-cache  
256 KB L2 cache  
8 MB L3 cache





## Great Reality #4:

There's more to performance than asymptotic complexity

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Must optimize at multiple levels: algorithm, data representations, procedures, and loops

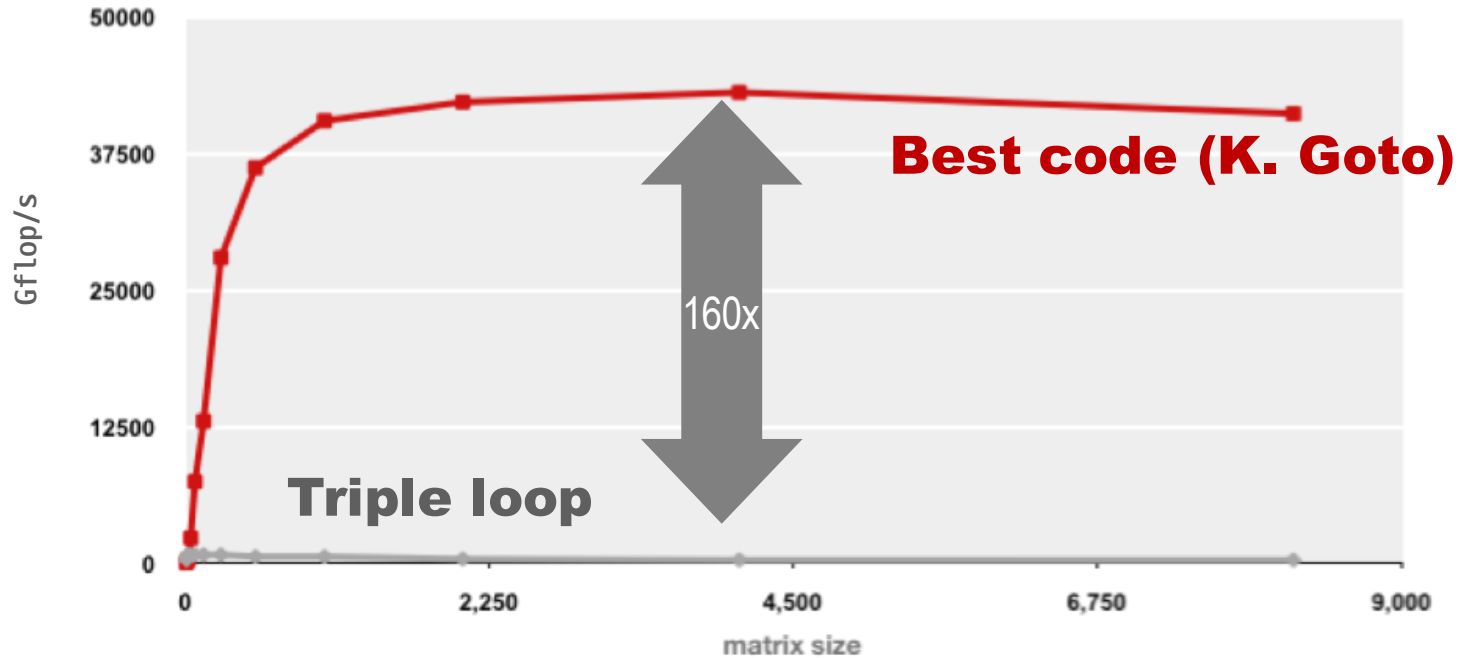
- Easily see 10:1 performance range depending on how code written

Must understand system to optimize performance

- How programs compiled and executed
- How to measure program performance and identify bottlenecks
- How to improve performance without destroying code modularity and generality

# Example Matrix Multiplication

Matrix-Matrix Multiplication (MMM) on 2 x Core 2 Duo 3 GHz (double precision)



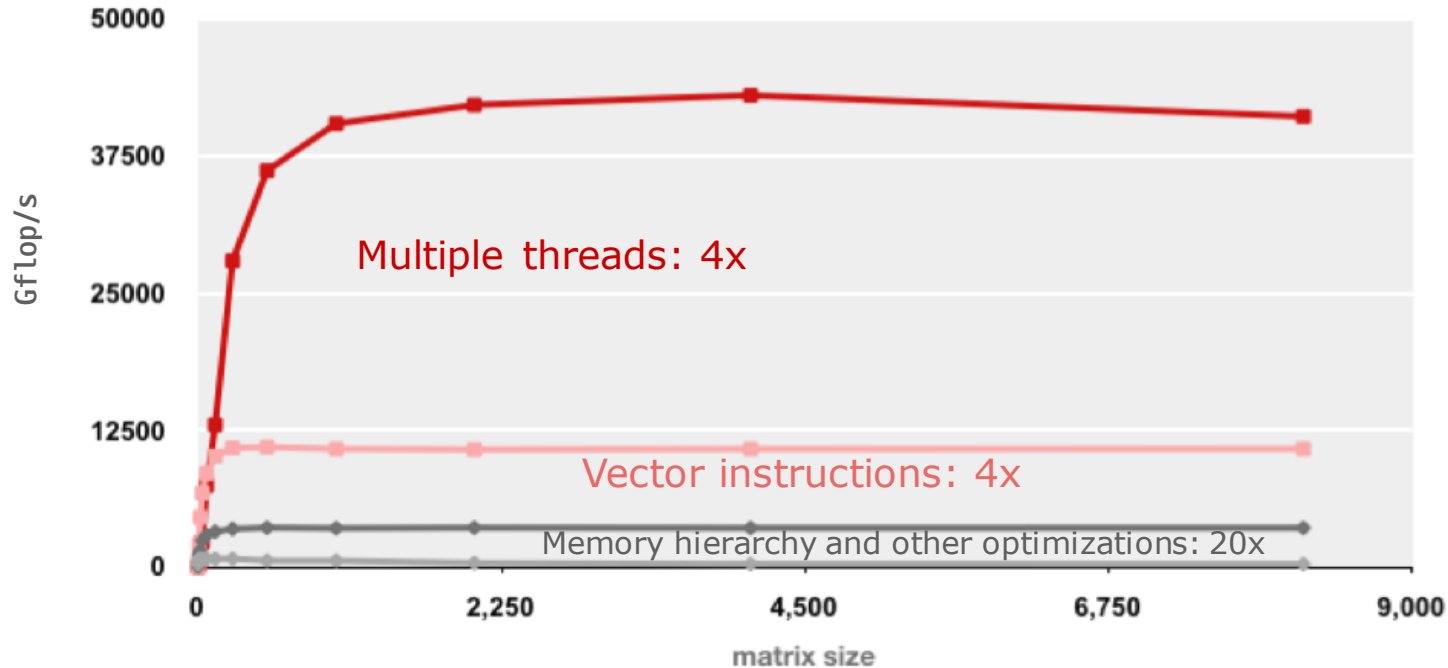
Standard desktop computer, vendor compiler, using optimization flags

- Both implementations have exactly the same operations count ( $2n^3$ )

What is going on?

# MMM Plot: Analysis

Matrix-Matrix Multiplication (MMM) on 2 x Core 2 Duo 3 GHz (double precision)



Reason for 20x:

- Blocking or tiling, loop unrolling, array scalarization, instruction scheduling, search to find best choice

Effect: fewer register spills, L1/L2 cache misses, and TLB misses

They need to get data in and out

- I/O system critical to program reliability and performance

They communicate with each other over networks

- Many system-level issues arise in presence of network
  - Concurrent operations by autonomous processes
  - Coping with unreliable media
  - Cross platform compatibility
  - Complex performance issues