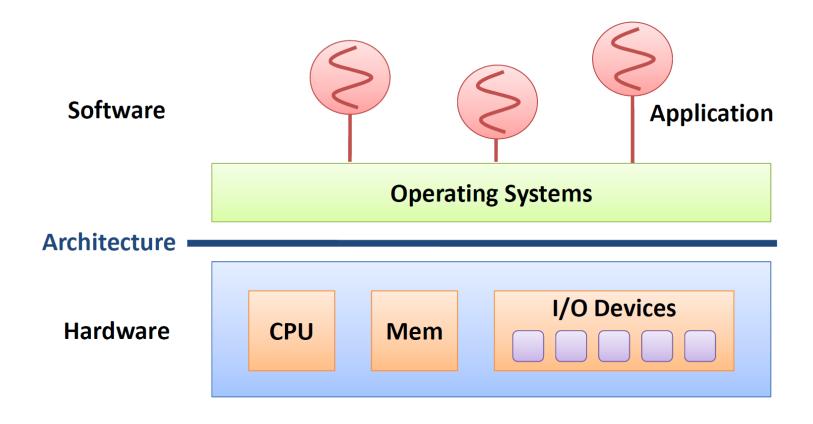
# OPERATING SYSTEM REVIEW (PROCESS)

Jo, Heeseung

# Operating system?

Computer systems internals



# **PROCESSES**

Jo, Heeseung

#### What Is The Process?

Program?

VS.

Process?

VS.

Processor?

VS.

Task? Job?

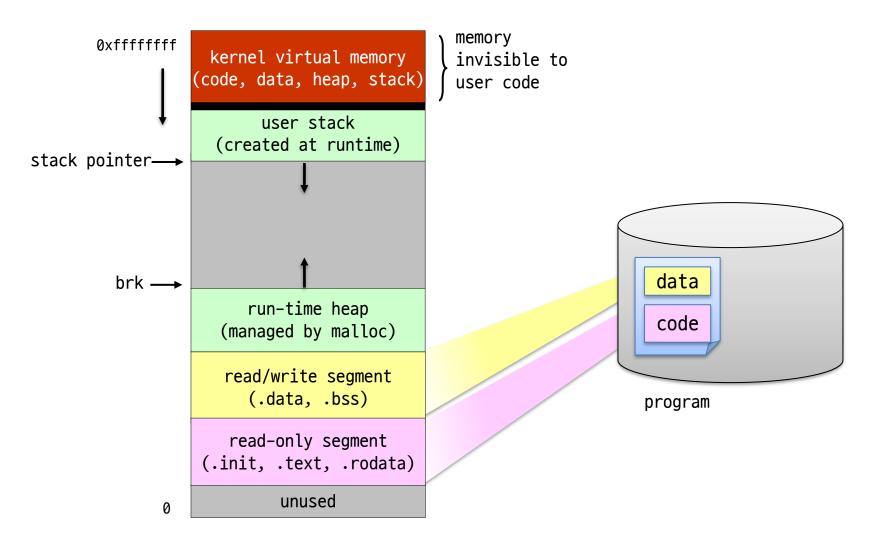
# Process Concept (1)

#### What is the process?

- An instance of a program in execution
- An encapsulation of the flow of control in a program
- A dynamic and active entity
- The basic unit of execution and scheduling
- A process is named using its process ID (PID)
- A process includes:
  - CPU contexts (registers)
  - OS resources (memory, open files, etc.)
  - Other information (PID, state, owner, etc.)

# Process Concept (2)

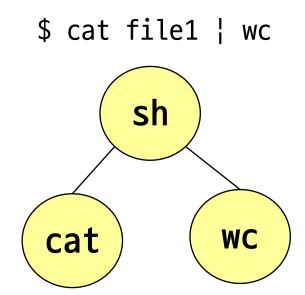
#### Process in memory



## Process Creation (1)

#### Process hierarchy

- One process can create another process: parent-child relationship
- UNIX calls the hierarchy a "process group"
- Windows has no concept of process hierarchy
- Browsing a list of processes:
  - ps in UNIX
  - taskmgr (Task Manager) in Windows



## Process Creation (2)

#### Process creation events

- Calling a system call
  - fork() in POSIX, CreateProcess() in Win32
  - Shells or GUIs use this system call internally
- System initialization
  - *init* process
  - PID 1 process

# Process Creation (3)

#### Resource sharing

- Parent may inherit all or a part of resources and privileges for its children
  - UNIX: User ID, open files, etc.

#### Execution

 Parent may either wait for it to finish, or it may continue in parallel

# cat wc

\$ cat file1 | wc

#### Address space

 Child duplicates the parent's address space or has a program loaded into it

#### **Process Termination**

#### Process termination events

- Normal exit (voluntary)
- Error exit (voluntary)
- Fatal error (involuntary)
  - Exceed allocated resources
  - Segmentation fault
  - Protection fault, etc.
- Killed by another process (involuntary)
  - By receiving a signal

```
#include <stdio.h>
int main()
{
    int i, fd;
    char buf[100];
    fd=open("a.txt", "r");
    if (fd==NULL)
        return -1;
    read(fd, buf, 1000);
    return 0;
```

# fork()

#### fork() system call

- Creating a child process
- Copy the whole virtual address space of parent to create a child process
- Copy internal data structures to manage a child process
- Parent get the pid of a child
- Child get 0 value

# fork()

```
#include <unistd.h>
                                   int main()
#include <sys/types.h>
                                       int pid;
#include <unistd.h>
                                       pid = fork();
                                       if (pid == 0)
int main()
                                         /* child */
                                         printf ("Child of %d is %d\n",
{
                                                 getppid(), getpid());
                                       else
      int pid;
                                         /* parent */
                                         printf ("I am %d. My child is %d\n",
                                                 getpid(), pid);
      pid = fork();
      if (pid == 0)
          /* child */
          printf ("Child of %d is %d\n",
                      getppid(), getpid());
      else
          /* parent */
          printf ("I am %d. My child is %d\n",
                      getpid(), pid);
```

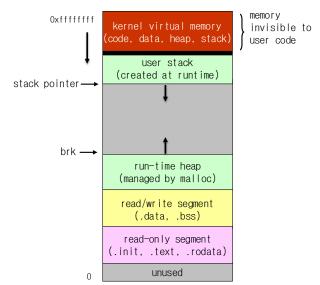
#include <sys/types.h>

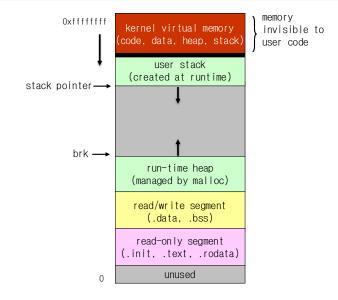
# fork(): Example Output

```
% ./a.out
I am 30000. My child is 30001.
Child of 30000 is 30001.

% ./a.out
Child of 30002 is 30003.
I am 30002. My child is 30003.
```

# fork() and Virtual Address Space





# Why fork()?

Very useful when the child ...

- Is cooperating with the parent
- Relies upon the parent's data to accomplish its task
- Example: Web server

```
While (1) {
    int sock = accept();
    if ((pid = fork()) == 0) {
        /* Handle client request */
    } else {
        /* Close socket */
    }
}
```

## Zombie vs. orphan process

#### Zombie process (defunct process)

- A process that completed execution (via the exit system call) but still has an entry in the process table
- This occurs for the child processes, where the entry is still needed to allow the parent process to read its child's exit status

```
int main() {
                                      [ijunseog-ui-MacBook-Pro: $ ./a.out &
                                      부모 PID : 60152, pid : 60153
    pid t childPid;
                                      자식 시작 PID : 60153
                                      ijunseog-ui-MacBook-Pro: shall hall jamana $ 자식 종료
                                      ps aux | grep 'Z'
    childPid = fork();
                                      USER
                                                      PID %CPU %MEM
                                                                       VSZ
                                                                                 TT STAT STARTED
                                                                                                     TIME COMMAND
                                                     60153
                                                           0.0 0.0
                                                                               0 s000 (Z)
                                                                                           7:16PM
                                                                                                   0:00.00 (a.out)
    if (childPid > 0) { // parent process
        printf("parent PID : %ld, pid : %d\n",(long)getpid(), childPid);
        sleep(30);
        printf("parent exit\n");
        exit(0);
    else if (childPid == 0){ // 자식 코드
        printf("child PID : %ld\n", (long)getpid());
        sleep(1);
        printf("child exit\n");
        exit(0);
    return 0;
}
```

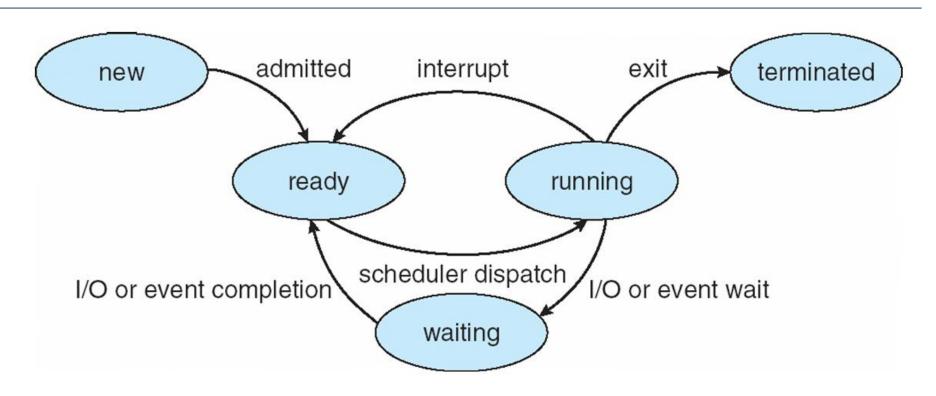
## Zombie vs. orphan process

#### Orphan process

- A process whose parent process has finished or terminated, though it remains running itself
- Any orphaned process will be immediately adopted by the special init system process

```
부모 PID: 46797, pid: 46798
                                                        자식 시작
                                                        자식 PID: 46798 부모 PID: 46797
int main() {
                                                        자식 PID: 46798 부모 PID: 46797
                                                        자식 PID: 46798 부모 PID: 46797
    pid t childPid;
                                                        ijunseog-ui-MacBook-Pro:s ------ <$ 자식 PID : 46798 부모 PID : 1
                                                        자식 PID : 46798 부모 PID : 1
    int i;
                                                        자식 PID : 46798 부모 PID : 1
                                                        자식 PID : 46798 부모 PID : 1
                                                        자식 PID : 46798 부모 PID : 1
    childPid = fork();
                                                        자식 PID : 46798 부모 PID : 1
                                                        자식 PID : 46798 부모 PID : 1
                                                        자 식 종 료
    if (childPid > 0) { // parent process
        printf("parent PID : %ld, pid : %d\n",(long)getpid(), childPid);
        sleep(2);
        printf("parent exit\n");
        exit(0);
    else if (childPid == 0){ // child process
        for(i=0;i<10;i++) {
             printf("child PID : %ld parent PID : %ld\n",(long)getpid(), (long)getppid());
             sleep(1);
        printf("child exit\n");
        exit(0);
```

# Process State Transition (1)



# **THREADS**

Jo, Heeseung

## Rethinking Processes

#### What's similar in these cooperating processes?

- They all use (share?) the same code and data (address space)
- They all use the same privilege
- They all use the same resources (files, sockets, etc.)

#### What's different?

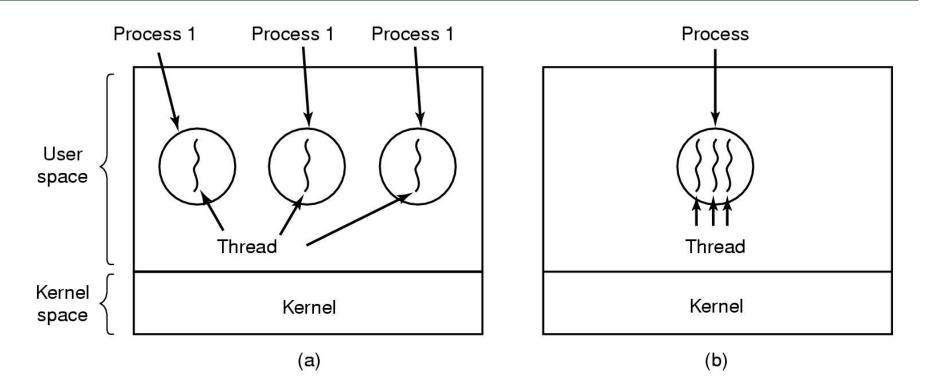
• Each has its own hardware execution state: PC, registers, SP, and stack

# Key Idea (1)

Separate the concept of a process from its execution state

- Process: address space, resources, other general process attributes
   e.g., privileges
- Execution state: PC, SP, registers, etc.
- This execution state is usually called
  - Thread
  - Lightweight process (LWP)
  - Thread of control

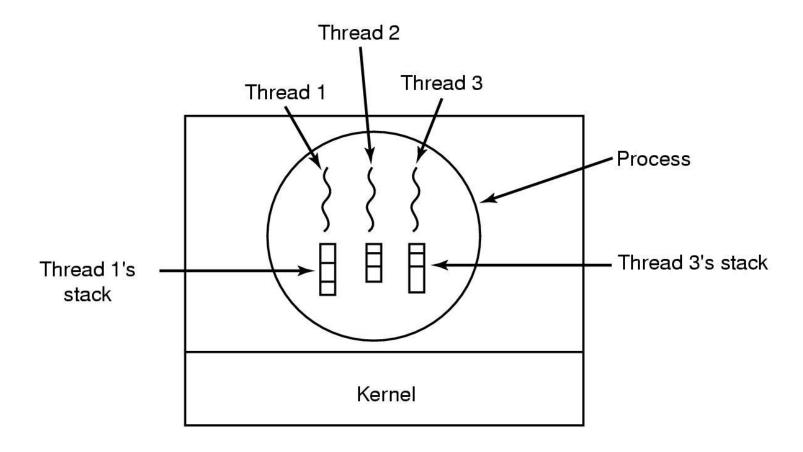
# Key Idea (2)



Per process items	Per thread items
Address space	Program counter
Global variables	Registers
Open files	Stack
Child processes	State
Pending alarms	
Signals and signal handlers	
Accounting information	

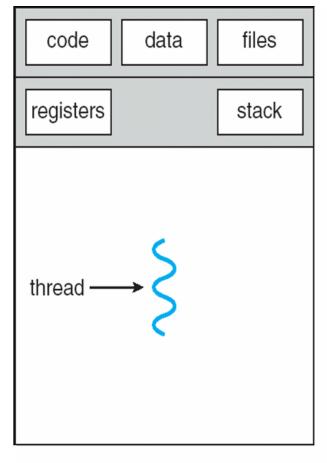
# Key Idea (3)

#### Each thread has its own stack

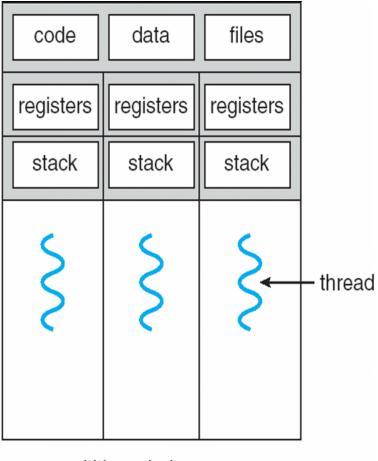


# Key Idea (4)

#### Each thread has its own stack



single-threaded process



multithreaded process

#### What is a Thread?

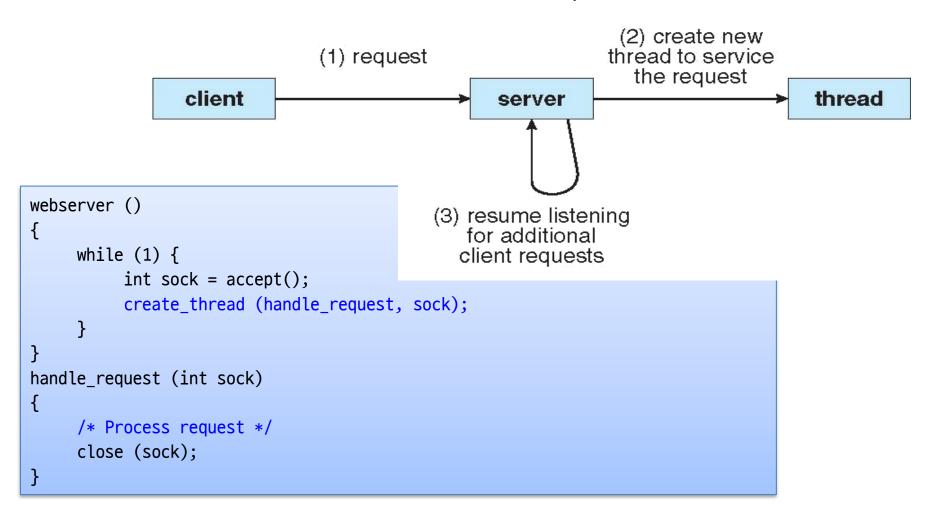
#### A thread of control (or a thread)

- A sequence of instructions being executed in a program
- Usually consists of
  - A program counter (PC), general registers
  - A stack to keep track of local variables and return addresses
- Threads share the process instructions and most of its data
  - A change in shared data by one thread can be seen by the other threads in the process
- Threads also share most of the OS state of a process

#### Concurrent Servers: Threads

#### Using threads

We can create a new thread for each request

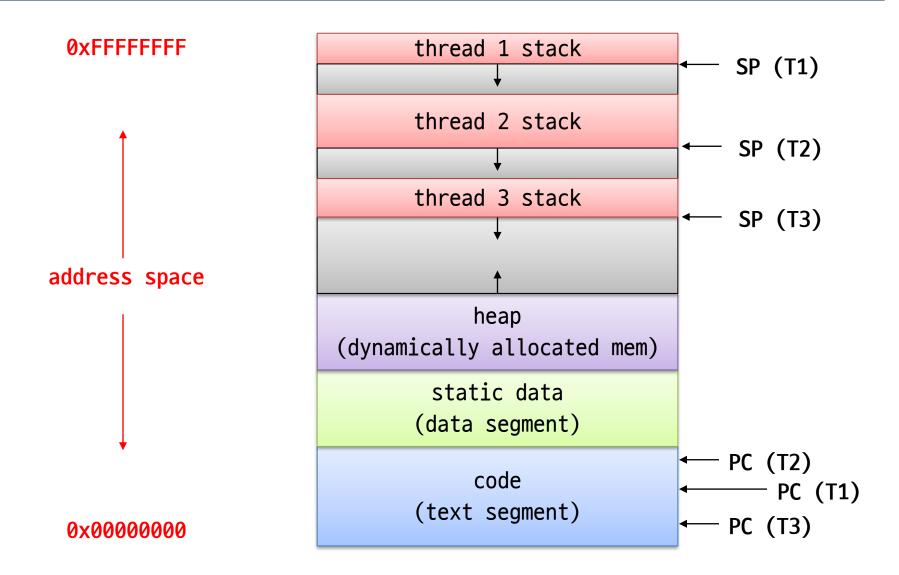


# Multithreading

#### Benefits

- Creating concurrency is cheap
  - Time and memory consumption
- Improves program structure
- Higher throughput
  - By overlapping computation with I/O operations
- Better responsiveness (User interface / Server)
  - Can handle concurrent events (e.g., web servers)
- Better resource sharing
- Utilization of multiprocessor architectures
  - Allows building parallel programs

# Address Space with Threads



# pthreads (1)

#### Thread creation/termination

# The Pthreads "hello, world" Program

```
#include <stdio.h>
                                             # gcc ex.c -lpthread
#include <pthread.h>
                                             # ./a.out
                                             main
void *threadfunc(void *vargp);
                                             Hello, world!
                                             main2
/* thread routine */
void *threadfunc(void *vargp) {
  sleep(1);
  printf("Hello, world!\n");
  return NULL;
int main() {
  pthread t tid;
  pthread_create(&tid, NULL, threadfunc, NULL);
  printf("main\n");
  pthread join(tid, NULL);
  printf("main2\n");
  sleep(2);
  return 0;
```

## Threading Issues (1)

```
fork() and exec() can be issue
```

When a thread calls fork()

- Does the new process duplicate all the threads?
- Is the new process single-threaded?

Some UNIX systems support two versions of fork()

- In pthreads,
  - fork() duplicates only a calling thread
- In the Unix international standard,
  - fork() duplicates all parent threads in the child
  - fork1() duplicates only a calling thread

Normally, exec() replaces the entire process

If a thread call exit()?

If the main thread dies(return, exit()) before child threads?