

COMPUTER ARCHITECTURE REVIEW

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What You Learned

How programs are translated into the machine language

- And how the hardware executes them

The hardware/software interface

What determines program performance

- And how it can be improved

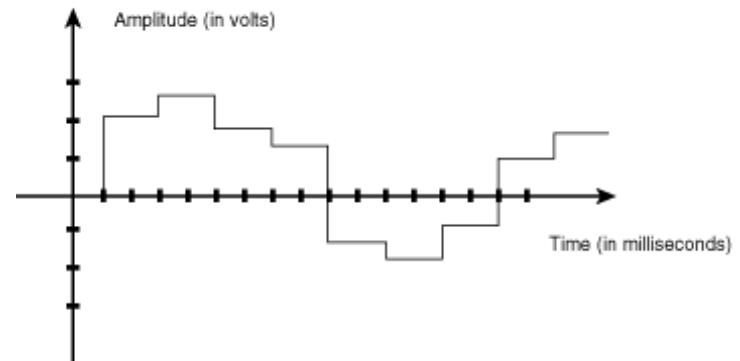
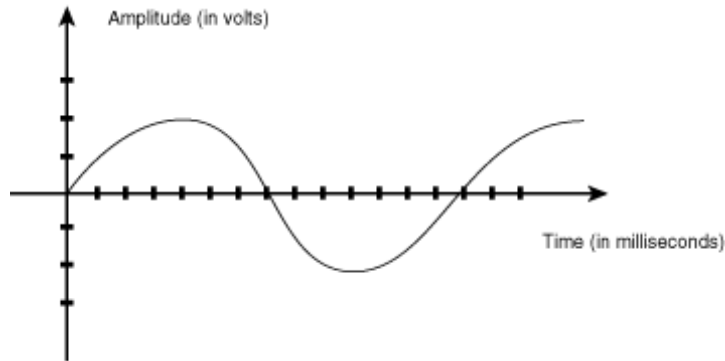
How hardware designers improve performance

What is parallel processing

Introduction

The advent of the digital age

- Analog vs. digital?



- Compact disc (CD)
 - 44.1 KHz, 16-bit, 2-channel
- MP3
 - A digital audio encoding with lossy data compression

Representing Information

Information = Bits + Context

- Computers manipulate representations of things
- Things are represented as binary digits
- What can you represent with N bits?
 - 2^N things
 - Numbers, characters, pixels, positions, source code, executable files, machine instructions, ...
 - Depends on what operations you do on them

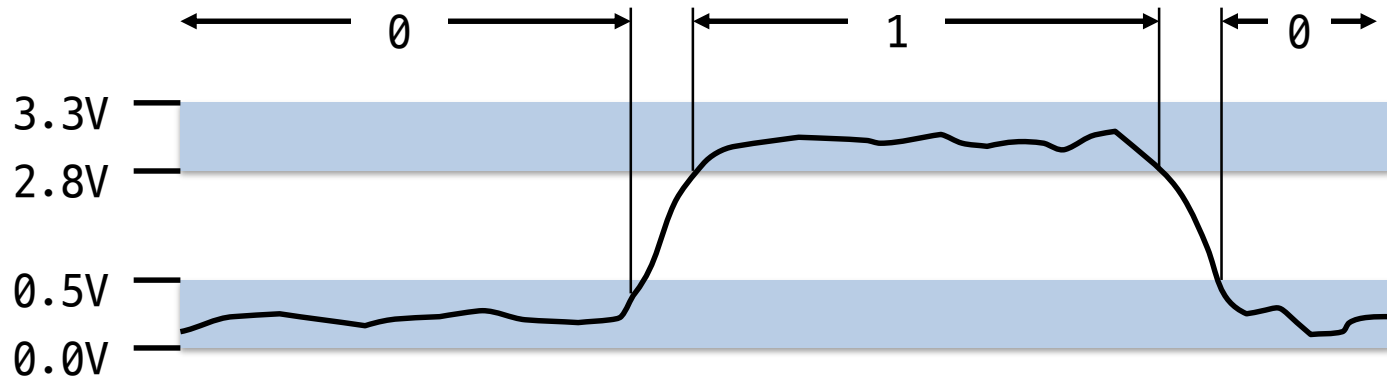
	01110011	01100101	01101101	01101001	01110011	01100101	01101101	01101001
(char)	's'	'e'	'm'	'i'	's'	'e'	'm'	'i'
(int)	1768777075				1768777075			
(double)	7.03168990329170808178... x 10 ¹⁹⁹							

Binary Representations

Why not base 10 representation?

- Easy to store with bistable elements
- Straightforward implementation of arithmetic functions
- Reliably transmitted on noisy and inaccurate wires

Electronic implementation



Encoding Byte Values

Byte = 8 bits

- Binary: 00000000_2 to 11111111_2
- Octal: 000_8 to 377_8
 - An integer constant that begins with 0 is an octal number in C
- Decimal: 0_{10} to 255_{10}
 - First digit must not be 0 in C
- Hexadecimal: 00_{16} to FF_{16}
 - Base 16 number representation
 - Use characters '0' to '9' and 'A' to 'F'
 - Write $FA1D37B_{16}$ in C as `0xFA1D37B` or `0xfa1d37b`

Hex	Decimal	Binary
0	0	0000
1	1	0001
2	2	0010
3	3	0011
4	4	0100
5	5	0101
6	6	0110
7	7	0111
8	8	1000
9	9	1001
A	10	1010
B	11	1011
C	12	1100
D	13	1101
E	14	1110
F	15	1111

Boolean Algebra (1)

Developed by George Boole in 1849

- Algebraic representation of logic
 - Encode "True" as 1 and "False" as 0

And

- $A \& B = 1$ when both $A=1$ and $B=1$

$\&$	0	1
0	0	0
1	0	1

Or

- $A | B = 1$ when either $A=1$ or $B=1$

$ $	0	1
0	0	1
1	1	1

Not

- $\sim A = 1$ when $A=0$

\sim	
0	1
1	0

Exclusive-Or (Xor)

- $A \wedge B = 1$ when either $A=1$ or $B=1$, but not both

\wedge	0	1
0	0	1
1	1	0

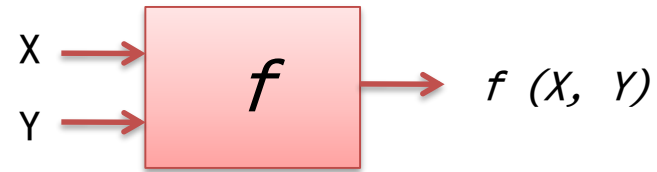
Boolean Algebra (2)

0	0	1	1
0	1	0	1

X
Y

0	0	0	0
0	0	0	1
0	0	1	0
0	0	1	1
0	1	0	0
0	1	0	1
0	1	1	0
0	1	1	1
1	0	0	0
1	0	0	1
1	0	1	0
1	0	1	1
1	1	0	0
1	1	0	1
1	1	1	0
1	1	1	1

Constant 0
 X & Y ; AND
 $\sim (X \rightarrow Y)$
 X
 $\sim (Y \rightarrow X)$
 Y
 X ^ Y ; XOR
 X | Y ; OR
 $\sim (X | Y)$; NOR
 $\sim (X \wedge Y)$; X-NOR
 $\sim Y$
 Y \rightarrow X
 $\sim X$
 X \rightarrow Y ; Implication
 $\sim (X \& Y)$; NAND
 Constant 1

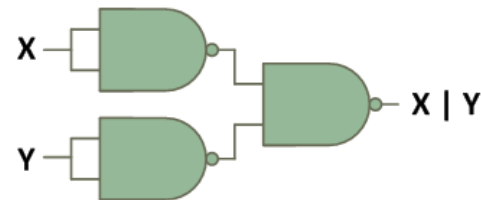
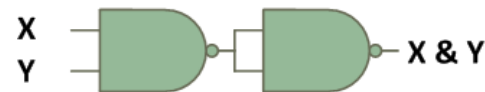


Basic operations: AND(&), OR(|), NOT(~)

$$X \wedge Y = (X \& \sim Y) | (\sim X \& Y)$$

$$X \rightarrow Y = \sim X | Y$$

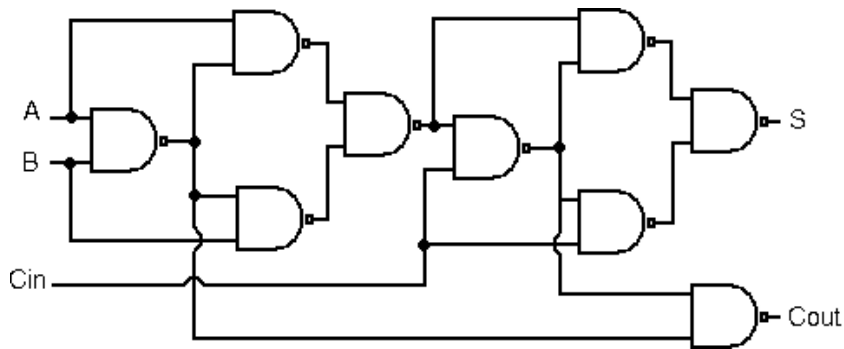
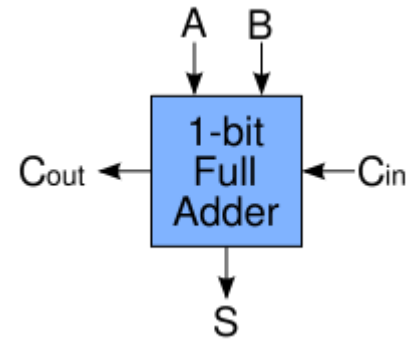
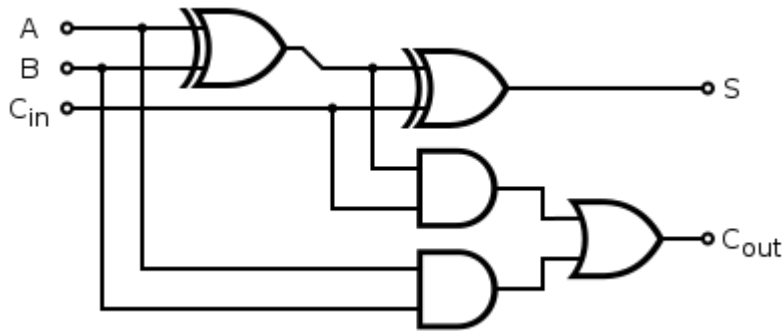
A complete set: NAND = $\sim (X \& Y)$



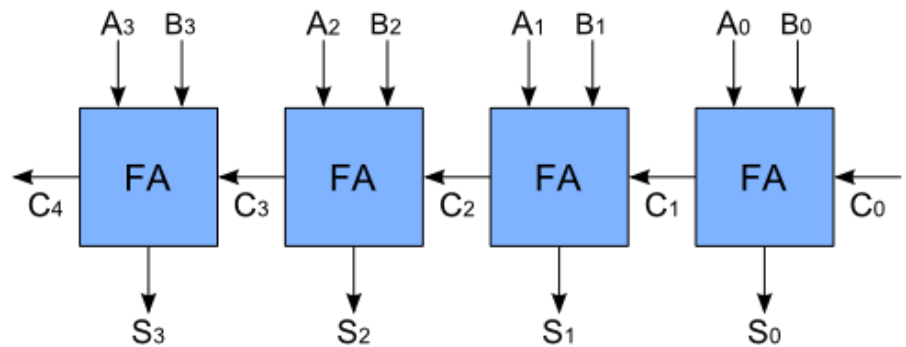
Combinational Logic

Adder

Full Adder



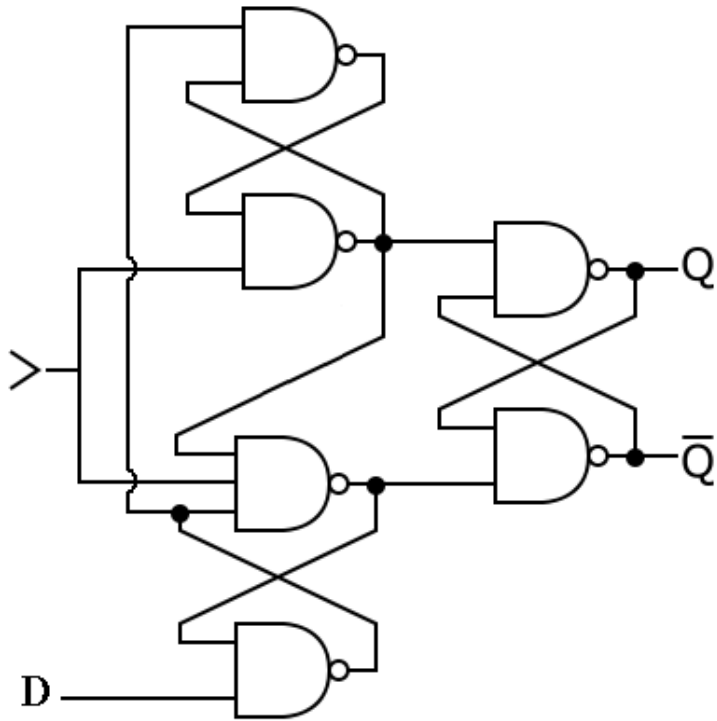
Full Adder (NAND version)



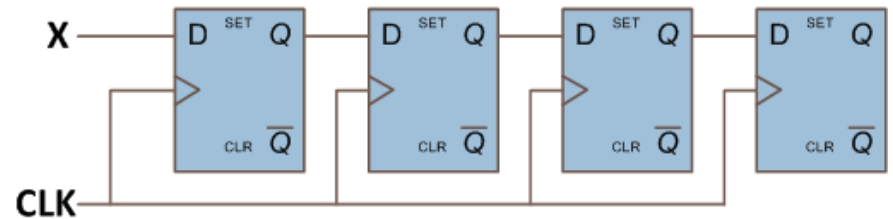
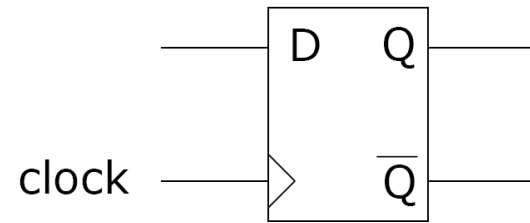
4-bit Ripple Carry Adder

Sequential Logic

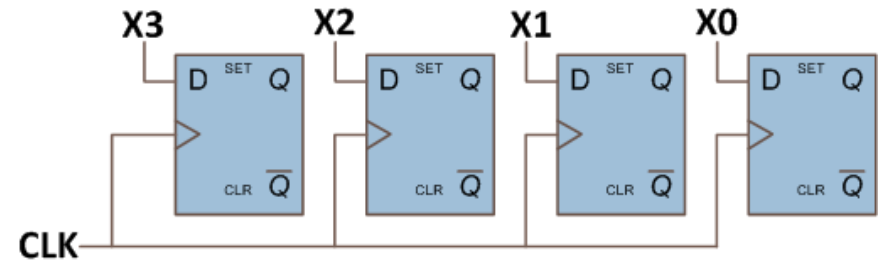
Flip-flops



Edge triggered D flip-flop



Shifter

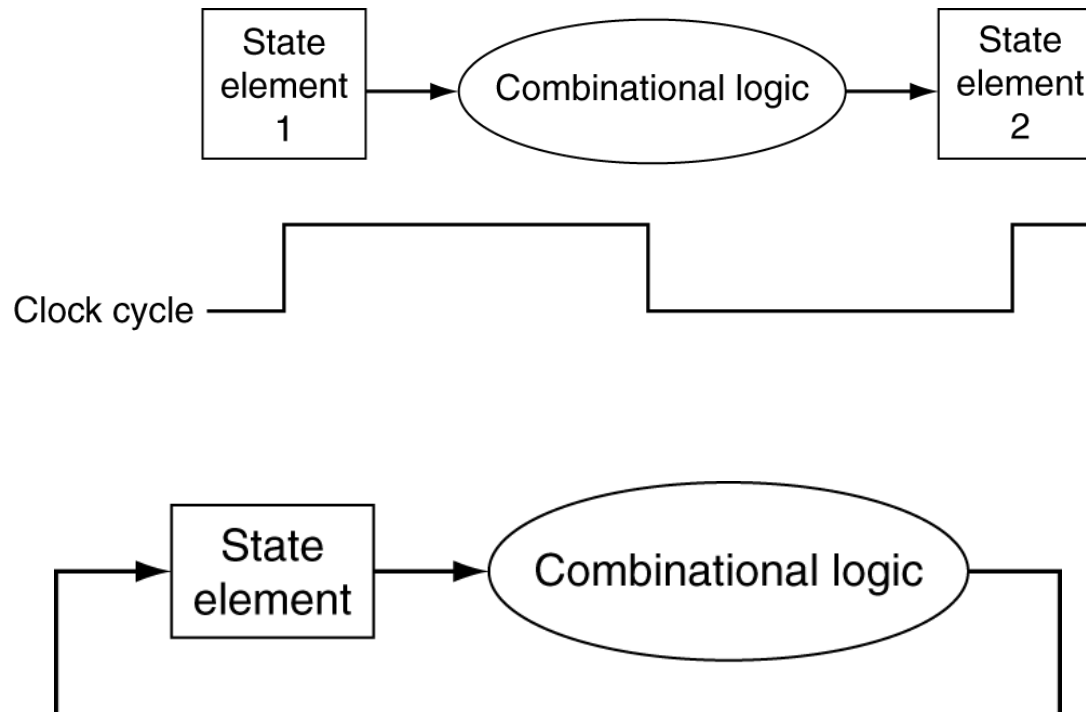


4-bit register

Clocking Methodology

Combinational logic transforms data during clock cycles

- Between clock edges
- Input from state elements, output to state element
- Longest delay determines clock period

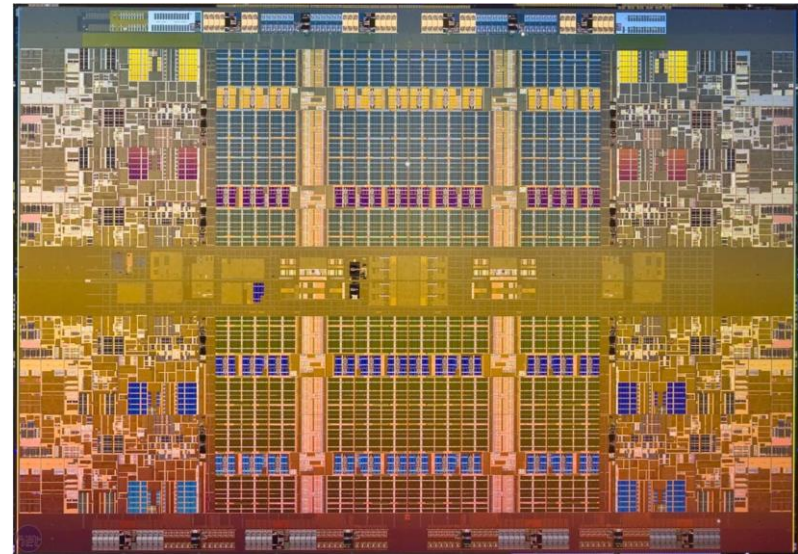


Digital Systems

Summary

- Boolean algebra is a mathematical foundation for modern digital systems
- Boolean algebra provides an effective means of describing circuits built with switches
 - Claude Shannon in the late 1930's
- You can build any digital systems with NAND gates
- A NAND gate can be easily built with CMOS transistors
- The transistor is the basic building block for digital systems

Intel Xeon 7560 (8-core): 2.3B transistors



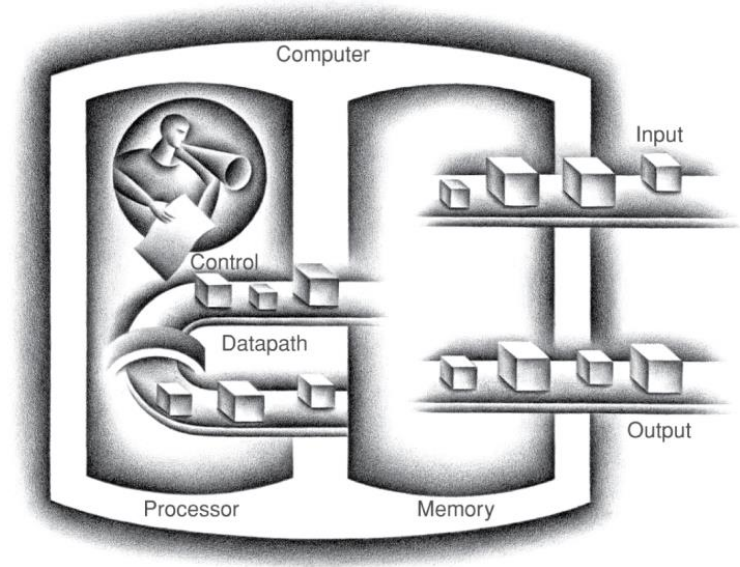
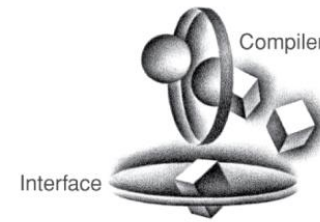
Components of a Computer

Same components for all kinds of computer

- Desktop, server, embedded

Input/output includes

- User-interface devices
 - Display, keyboard, mouse
- Storage devices
 - Hard disk, CD/DVD, flash
- Network adapters
 - For communicating with other computers



Understanding Performance

Algorithm

- Determines number of operations executed

Programming language, compiler, architecture

- Determine number of machine instructions executed per operation

Processor and memory system

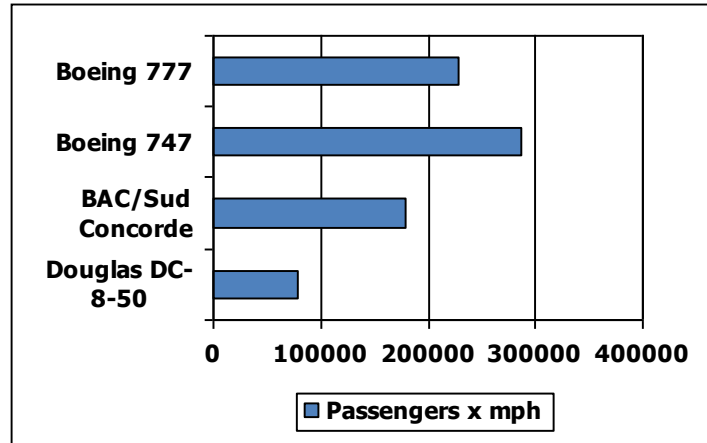
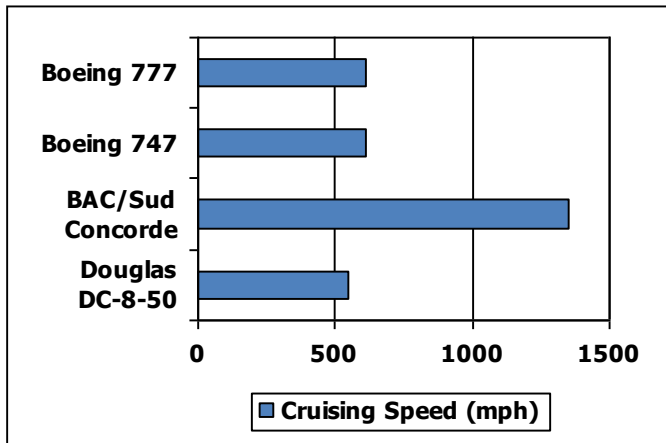
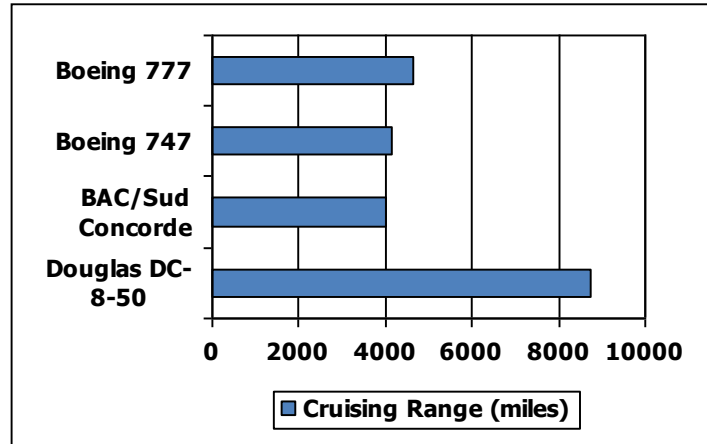
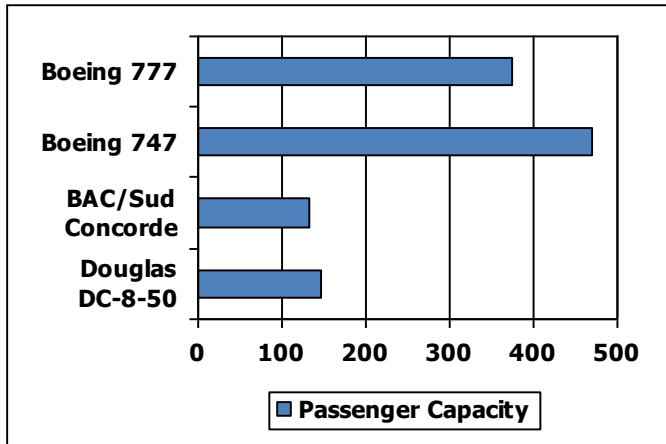
- Determine how fast instructions are executed

I/O system (including OS)

- Determines how fast I/O operations are executed

Defining Performance

Which airplane has the best performance?



Response Time and Throughput

Response time

- How long it takes to do a task

Throughput

- Total work done per unit time
 - e.g., tasks/transactions/... per hour

How are response time and throughput affected by

- Replacing the processor with a faster version?
- Adding more processors?

Levels of Program Code

High-level language

- Level of abstraction closer to problem domain
- Provides for productivity and portability

Assembly language

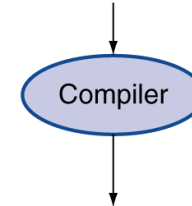
- Textual representation of instructions

Hardware representation

- Binary digits (bits)
- Encoded instructions and data

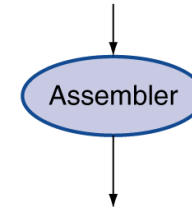
High-level
language
program
(in C)

```
swap(int v[], int k)
{int temp;
  temp = v[k];
  v[k] = v[k+1];
  v[k+1] = temp;
}
```



Assembly
language
program
(for MIPS)

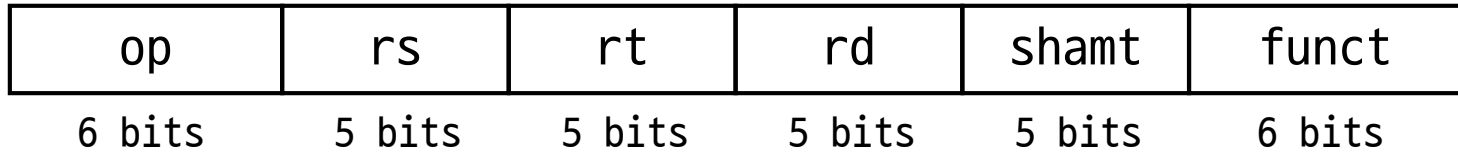
```
swap:
  muli $2, $5,4
  add $2, $4,$2
  lw $15, 0($2)
  lw $16, 4($2)
  sw $16, 0($2)
  sw $15, 4($2)
  jr $31
```



Binary machine
language
program
(for MIPS)

```
000000001010000100000000000011000
000000000000110000001100000100001
100011000110001000000000000000000
100011001111001000000000000000100
101011001111001000000000000000000
101011000110001000000000000000100
00000011111000000000000000001000
```

MIPS R-format Instructions



Instruction fields

- op: operation code (opcode)
- rs: first source register number
- rt: second source register number
- rd: destination register number
- shamt: shift amount (00000 for now)
- funct: function code (extends opcode)

R-format Example

op	rs	rt	rd	shamt	funct
6 bits	5 bits	5 bits	5 bits	5 bits	6 bits

add \$t0, \$s1, \$s2

special	\$s1	\$s2	\$t0	0	add
---------	------	------	------	---	-----

0	17	18	8	0	32
---	----	----	---	---	----

000000	10001	10010	01000	00000	100000
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$$00000010001100100100000000100000_2 = 02324020_{16}$$

Translation and Startup

